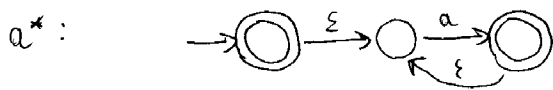
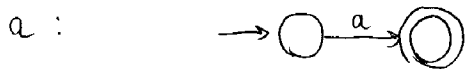


A SIMILAR EXAMPLE TO PROBLEM 1 :

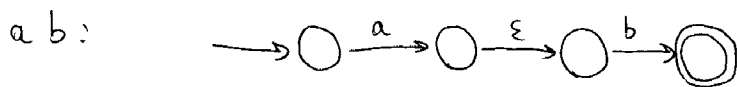
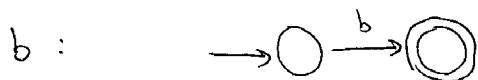
(1)

$$L = a^* U (a b)^*$$

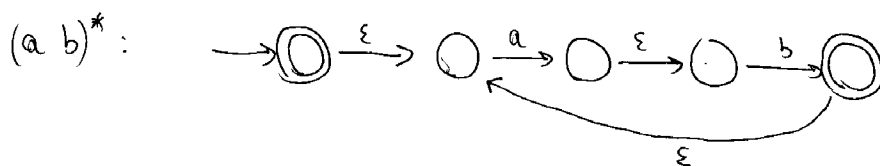
FOLLOWING THE STEPS OF LEMMA 1.55 WE GET



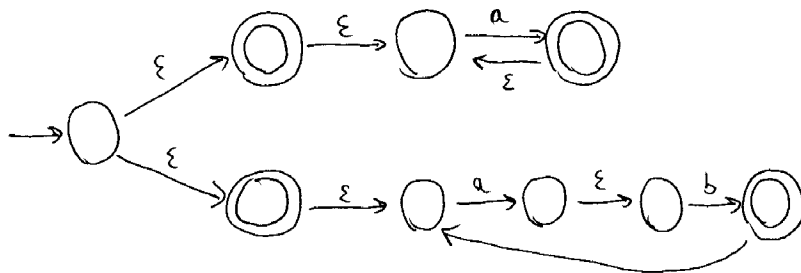
(SEE PAGE 62 FOR THIS)



(SEE PAGE 61)



a\* U (a b)\* :

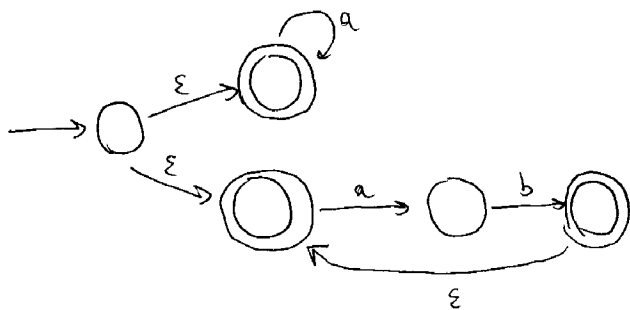


(SEE PAGE 59)

THIS IS OBTAINED FOLLOWING EVERY SINGLE STEP PRECISELY

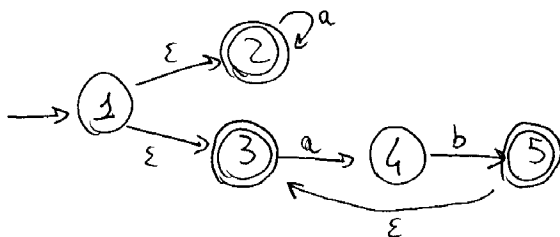
BUT IF YOU SEE A CLEAR SIMPLIFICATION, YOU CAN SIMPLIFY TO REDUCE THE NUMBER OF STATES AND MAKE THE PROBLEM EASIER FOR THE NEXT STEP

FOR EXAMPLE



THE 2<sup>ND</sup> PART REQUIRES TO TRANSFORM THIS DFA

(2)



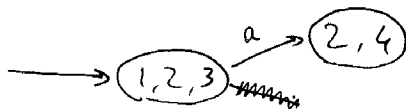
INTO A DFA

WE START FROM 1, THE START STATE

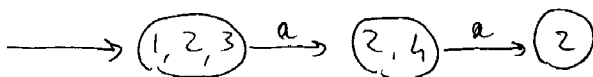
A "ε" ALLOWS TO MOVE DIRECTLY TO THE DESTINATION STATES, SO BEING IN 1 IS LIKE BEING IN 2 AND 3 AT THE SAME TIME. FOR THIS REASON WE GET A COMBINED STATE



TO UNDERSTAND HOW IT WORKS, PUT YOUR FINGERS ON 1, 2 AND 3. WHERE CAN YOU GO WITH AN "a"? YOU END UP IN 2 AND IN 4. THIS MEANS THAT THERE IS A STATE (2, 4):

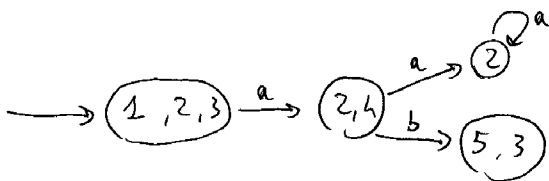


NOW WE ARE DONE WITH (1, 2, 3). WHERE DO WE GO FROM (2, 4) READING AN "a"? ONLY TO 2, SO THIS GIVES A NEW STATE WITH 2 ONLY

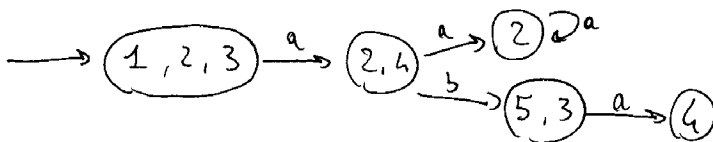


FROM 2, WITH a WE LOOP AROUND.

FROM (2, 4) WITH b WE ARRIVE TO 5. BUT 5 HAS AN OUTGOING ε, WHICH MEANS THAT BEING IN 5 IS LIKE BEING IN 3 TOO

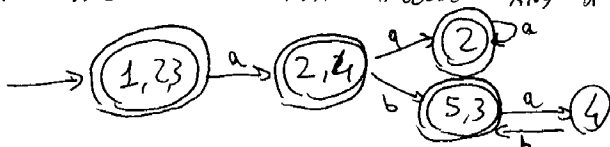


NOW FROM 3 WE CAN GO TO 4 WITH AN a, BUT THIS IS NOT (2, 4), SO



FINALLY FROM 4 WE CAN GO BACK TO 5, 3 WITH A b.

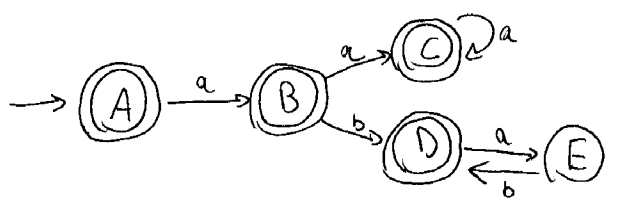
THE FINAL STATES ARE THOSE THAT INCLUDE ANY OF THE ORIGINAL ONES:



IT IS VERY IMPORTANT TO DOUBLE CHECK THE RESULT, BECAUSE IT IS VERY EASY TO FORGET SOMETHING...

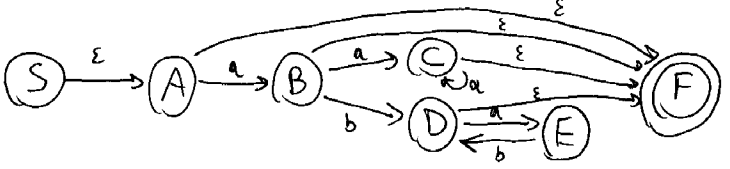
THE FINAL PART CONSISTS IN CONVERTING BACK THE DFA

3

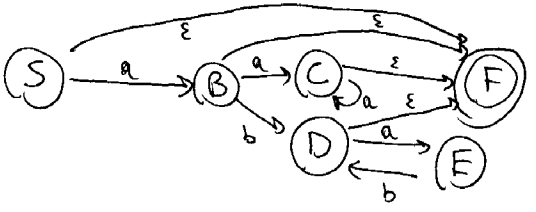


INTO A REGULAR EXPRESSION.

WE NEED TO ADD A START STATE AND A FINAL STATE WHICH WE SHOULD BE CONNECTED TO ALL FINAL STATES WITH AN  $\epsilon$ -ARROW

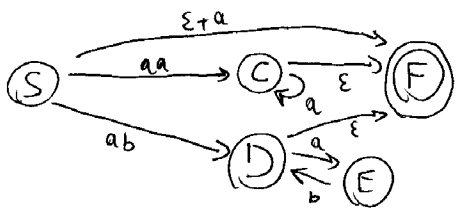


WE ELIMINATE ONE STATE AT A TIME STARTING FROM A (THE ORDER DOES NOT MATTER) IF A IS NOT THERE, WE STILL NEED TO REACH THE OTHER STATES IN THE SAME WAY AS BEFORE, SO " $\epsilon a$ " WOULD BRING TO B AND " $\epsilon \epsilon$ " TO F

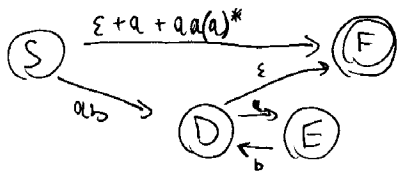


NOW WE REMOVE B. LET'S LOOK AT ALL THE DESTINATIONS OF B:

- BF: THIS MEANS THAT SF SHOULD INCLUDE " $a \epsilon$ "
- BC: " " " SC " " " $aa$ "
- SD: " " " SD " " " $ab$ "

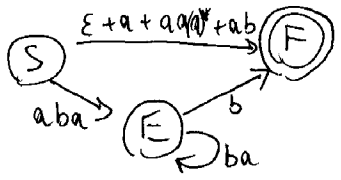


NOW WE REMOVE C: THE LOOP TRANSLATES INTO  $a^*$  SO WE GET



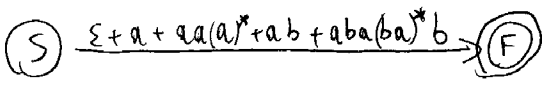
MAYBE REMOVING E WOULD MAKE THE LIFE EASIER, BUT JUST TO MAKE THIS EXAMPLE MORE INTERESTING WE FOLLOW THE DIFFICULT PATH...

WHEN WE REMOVE D, WE SHOULD PAY ATTENTION ON THE THE PATH  $S \rightarrow D \rightarrow F$  AND ALSO ON SDE AND EDE:

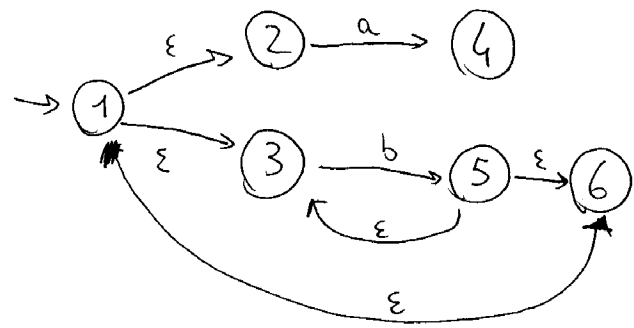


THIS WAS THE MOST CRITICAL STEP, PLEASE MAKE SURE YOU UNDERSTAND ALL DETAILS

FINALLY:



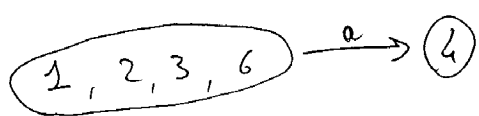
ONE MORE EXAMPLE FROM NFA TO DFA



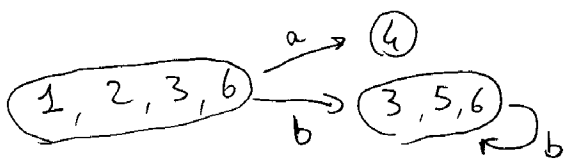
AT THE BEGINNING WE ARE IN 1, BUT BECAUSE OF  $\epsilon$  WE ARE ALSO IN 2, 3 AND 6, SO THE FIRST STATE IS

1, 2, 3, 6

FROM 2 WITH AN a WE GO TO 4, SO



AND FROM 3, WE GO TO 5 WITH A b, BUT 5 IS ALSO 3 AND 6 BECAUSE OF THE  $\epsilon$ , SO

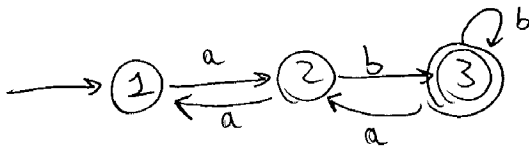


THIS SELF LOOP IS BECAUSE FROM 3 WE CAN AGAIN GO TO 3, 5, 6

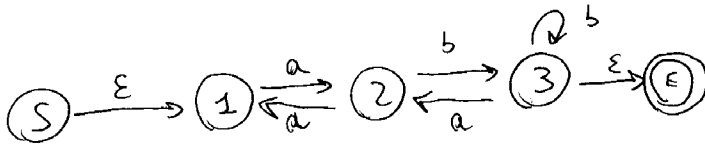
NOTHING ELSE TO DO

ONE MORE EXAMPLE FROM DFA TO REGULAR EXPRESSION

(5)

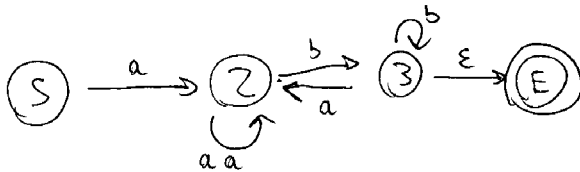


FIRST WE ADD A START AND END STATE



NOW WE ELIMINATE ONE STATE AT A TIME

IF WE ELIMINATE STATE 1, WE SHOULD STILL BE ABLE TO REACH 2 WITH AN  $a$  AND COME BACK TO 2 WITH TWO  $a$ 's (" $aa$ ")



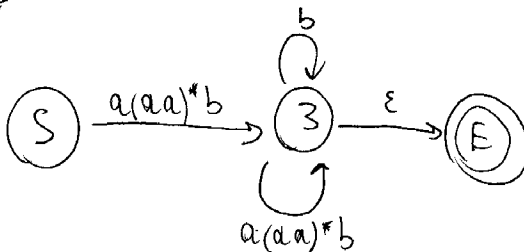
WHAT HAPPENS IF WE ELIMINATE 2?

WE SHOULD STILL BE ABLE TO REACH 3 WITH  $a(aa)^*b$

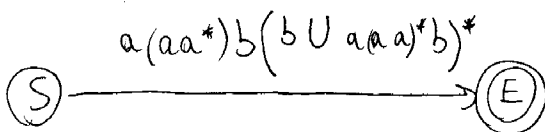
BUT WE ALSO NEED TO BE ABLE TO LOOP AROUND 3  $a(aa)^*b$

(BECAUSE FROM 3 WE CAN GO TO 2 WITH AN  $a$ , LOOP AROUND WITH  $aa$  AND COME BACK WITH  $b$ )

SO WE GET



FINALLY WE ELIMINATE 3. THERE ARE TWO WAYS TO LOOP AROUND 3, SO THAT REPRESENTS A UNION OPERATION



THEN THE ANSWER IS  $a(aa)^*b(b \cup a(aa)^*b)^*$