

RECCURENCE → CLOSED FORMULA → SUM <sup>(7)</sup>

R

CF

$S_n$

TOWER of HANOI (revisited)

GOAL

**R**  $T_0 = 0$

$T_n = 2T_{n-1} + 1$

**CF**

QUESTION

$T_n = 2^n - 1$

Divide by  $2^n$

$T_0/2^0 = 0$

$T_n/2^n = \frac{T_{n-1}}{2^{n-1}} + \frac{1}{2^n}$

$S_n = T_n/2^n$

R = RS

USE **CF**

**RS**  $S_0 = 0$

$S_n = S_{n-1} + \frac{1}{2^n}$   $n \geq 1$

$S_n = \frac{2^n - 1}{2^n}$

$S_n = \sum_{k=0}^n \frac{1}{2^k}$

$= 1 - \frac{1}{2^{n+1}}$  SUM

$S_n = \frac{a_1(1-q^{n+1})}{1-q}$  ;  $a_1 = \frac{1}{2}$  ;  $q = \frac{1}{2}$

$\sum_{k=0}^n \frac{1}{2^k} = 1 - \frac{1}{2^{n+1}}$

# GENERAL TECHNIQUE

for finding a CLOSED FORM  
for any RECURRENCE of a TYPE:  
with some  $n=0$  INITIAL CONDITION.

$$\textcircled{R} \quad a_n T_n = b_n T_{n-1} + C_n \quad \textcircled{n \geq 1}$$

BY REDUCING IT TO A CERTAIN SUM  
where

$a_n, b_n, C_n$  any sequences,  $n \geq 1$

## IDEA

Multiply R by a <sup>small</sup> SUMMATION FACTOR  $S_n, n \geq 1$

(we don't know yet what this factor is, but we will find it out)

$S_n = s(n) \quad s: \mathbb{N} \rightarrow \mathbb{R}$   
sequence

our (R)

$$a_n T_n = b_n T_{n-1} + C_n$$

$n \geq 1$

$S_n$

SUMMATION FACTOR

$$S_n a_n T_n = S_n b_n T_{n-1} + S_n C_n$$

\*

We WANT  $S_n$  to HAVE A PROPERTY

(GOAL):

GET RECURSIVE FORMULA

IT MEANS I TAKE a segue  $S_n$  inside that (P) holds

(P)

$$S_n b_n = S_{n-1} a_{n-1}$$

WE WRITE

(S)

$$S_n = S_n a_n T_n$$

$S_0 = S_0 a_0 T_0$   
(as in Tower of Hanoi)  
 $S_n = \frac{1}{2^n} T_n$

and RE-WRITE \* as (use (P))

(RS)

$$S_n = S_{n-1} + S_n C_n$$

$n \geq 1$

$$S_n = S_n b_n T_{n-1} + S_n C_n \quad \text{and} \quad S_n b_n = S_{n-1} a_{n-1} \quad (P)$$

$$S_m = s_0 a_0 T_0 + \sum_{k=1}^m S_k C_k$$

$$s_{n-1} a_{n-1} = s_n b_n$$

use  $P$

$$S_m = s_1 b_1 T_0 + \sum_{k=1}^m S_k C_k$$

By  $S$   $\downarrow$  det  
 $S_n = a_n s_n T_n$   
 we get

$$T_m = \frac{S_m}{a_m S_m}$$

We get  $T_m = \frac{1}{a_m S_m} \cdot S_m$

$$T_m = \frac{1}{S_m a_m} \left( s_1 b_1 T_0 + \sum_{k=1}^m S_k C_k \right)$$

$S$  closed

"SUM" closed formula for  $T_m$

NEXT STEP: FIND  $S_m$  in terms of  $a_m, b_m, C_m$

Q: HOW TO FIND  $S_n$  FACTOR.

Evaluate P for  $n \geq 1$

$$S_n = \frac{S_{n-1} a_{n-1}}{b_n}$$

$a_n, b_n$  are GIVEN

$$S_2 = \frac{S_1 a_1}{b_2} = S_1 \frac{a_1}{b_2}$$

$$S_3 = \frac{S_2 \cdot a_2}{b_3} = S_1 \frac{a_1 a_2}{b_2 b_3}$$

$$S_4 = \frac{S_3 a_3}{b_4} = S_1 \frac{a_1 a_2 a_3}{b_2 b_3 b_4}$$

BY MATH INDUCTION

$$S_n = S_1 \frac{a_1 a_2 \dots a_{n-1}}{b_2 b_3 b_4 \dots b_n}$$

CONSTANT summation FACTOR

R

$T_0$  - given initial cond.

$$a_n T_n = b_n T_{n-1} + C_n$$

$n \geq 1$

$a_n, b_n, C_n$   
for  $n \geq 1$

CF

$$T_n = \frac{1}{S_n a_n} \left( S_1 b_1 T_0 + \sum_{k=1}^n S_k C_k \right)$$

SUMMATION FACTOR

$$S_n = \frac{a_1 a_2 \dots a_{n-1}}{b_2 b_3 \dots b_n}$$

CONST =  $S_1$

Example Tower of Hanoi

$T_0 = 0$   
 $T_n = 2T_{n-1} + 1$

is a case where

$$a_n = 1, b_n = 2, C_n = 1$$

evaluate

$S_1 = \frac{1}{2}$

$$S_n = \frac{1}{\underbrace{2 \dots 2}_{n-1}} \cdot \underbrace{\left(\frac{1}{2}\right)}_{S_1} = \frac{1}{2^n}$$

$S_n = 2^{-n}$

Check

$$T_n = \frac{1}{a_n S_n} \left( s_i b_i T_0 + \sum_{k=1}^n S_k C_k \right)$$

for  $S_n = 2^{-n}$ ,  $a_n = 1$ ,  $b_n = 2$ ,  $C_n = 1$

all  $n$

$$T_0 = 0$$

$$T_n = \frac{1}{2^{-n}} \left( 0 + \sum_{k=1}^n \frac{1}{2^k} \right) \quad \text{closed SUM FORMULA}$$

$$T_n = 2^n \left( 1 - \frac{1}{2^n} \right)$$

$$T_n = 2^n - 1$$

→ geometric

$$S_n = \frac{a_0 (q^{n+1} - 1)}{q - 1}$$

$$S_n = \frac{a_0 (1 - q^{n+1})}{1 - q}$$

$$q = \frac{1}{2}$$

geometric

Closed

FORMULA

$$\sum_{k=1}^n \frac{1}{2^k} = 1 - \frac{1}{2^n}$$

# QUICKSORT

HOARE 1962

The average number of comparison stops made by the QS when applied to  $n$  items in RANDOM order is

$$C_0 = 0$$

$$C_n = (n+1) + \frac{2}{n} \sum_{k=0}^{n-1} C_k$$

$$C_1 = 2, C_2 = 5, C_3 = \frac{26}{3} \text{ etc..}$$

STEP 1

GET RID OF  $\Sigma$

in the recurrence

STEP 2

Find  $C_F$

(summation at least)

USE PREVIOUS TECHNIQUE!



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$$C_n = n+1 + \frac{2}{n} \sum_{k=0}^{n-1} C_k$$

$$n C_n = n^2 + n + 2 \sum_{k=0}^{n-1} C_k$$

$n > 1$

$$n C_n = n^2 + n + 2 \left( \sum_{k=0}^{n-2} C_k + C_{n-1} \right)$$

$$n C_n = n^2 + n + 2 \cdot \sum_{k=0}^{n-2} C_k + 2 \cdot C_{n-1}$$

Re-write  $\textcircled{0}$  for  $n = n-1$  we need

$$(n-1)C_{n-1} = (n-1)^2 + n-1 + 2 \sum_{k=0}^{n-2} C_k$$

$n^2 - 2n + 1 + n - 1$

$$(n-1)C_{n-1} = n^2 - n + 2 \sum_{k=0}^{n-2} C_k$$

SUBSTRACT  $\textcircled{2}$  from  $\textcircled{1}$

$$n C_n - (n-1)C_{n-1} = 2n + 2C_{n-1}$$

no  $\Sigma$ !

re-write it

(1/a)

We have:

→ we got a formula

(1)

$$S_n = \frac{T_n}{2^n}$$

$$S_n = S_{n-1} + \frac{1}{2^n}$$

summation

and

(2)

$$S_n = \frac{2^n - 1}{2^n}$$

by evaluating  
(RS)

(1) + (2)

$$T_n = 2^n \cdot S_n = \frac{2^n (2^n - 1)}{2^n} = 2^n - 1$$

$$T_n = 2^n - 1$$

end.

$$nC_n = (n-1)C_{n-1} + 2n + 2C_{n-1}$$

NEW 80  
FORMULA  
 $C_0 = 0$

$$= nC_{n-1} - C_{n-1} + 2n + 2C_{n-1}$$

$$= 2n + nC_{n-1} + C_{n-1} \quad \text{FORMULA}$$

$$nC_n = (n+1)C_{n-1} + 2n, \quad C_0 = 0$$

THIS IS

$$a_n T_n = b_n T_{n-1} + C_n$$

GENERAL  
RECURSION  
SOLVED  
BEFORE!

FOR

$$a_n = n, \quad b_n = n+1, \quad C_n = 2n$$

SUMMATION FACTOR

$$S_n = \frac{a_1 a_2 \dots a_{n-1}}{b_2 \dots b_n}$$

multiplied  
by a constant  
 $S_1$

$$S_n = \frac{1 \cdot 2 \dots (n-1)}{2 \dots (n-1) n (n+1)}$$

$$= \frac{2}{n(n+1)} \cdot S_1$$

$b_2 = 3$       $S_1 = \frac{2}{1 \cdot 2} = 1$       $S_1 = 1$

USE

$$T_n = \frac{1}{s_n a_n} \left( s, b, T_0 + \sum_{k=1}^n s_k c_k \right)$$

$$T_0 = c_0 = 0$$

use this for  $c_n$

$$a_n = n$$

$$b_n = n+1$$

$$c_n = 2n$$

$$c_n = \frac{1}{s_n \cdot n} \left( 0 + \sum_{k=1}^n 2k \cdot s_k \right)$$

$$s_k = \frac{2}{k(k+1)}$$

$$c_0 = 0$$

$$c_n = \frac{n(n+1)}{2n} \sum_{k=1}^n \frac{4k}{k(k+1)}$$

$$c_n = 2(n+1) \sum_{k=1}^n \frac{1}{k+1}$$

HARMONIC NUMBER  $H_n$

$$H_n = 1 + \frac{1}{2} + \dots + \frac{1}{n} = \sum_{k=1}^n \frac{1}{k}$$

**NAME:**  $k$ -th harmonic produced by a violin string is the fundamental tone produced by a string that is  $\frac{1}{k}$  times as long

$$\sum_{k=1}^n \frac{1}{k+1} = \sum_{1 \leq k \leq n} \frac{1}{k+1} = \sum_{1 \leq k-1 \leq n} \frac{1}{k}$$

put  
 $k=k+1$

$$= \sum_{2 \leq k \leq n+1} \frac{1}{k} = \sum_{k=1}^n \frac{1}{k} - \frac{1}{1} + \frac{1}{n+1}$$

$$\sum_{k=1}^n \frac{1}{k+1} = H_n - \frac{n}{n+1}$$

$$\frac{-n-1+1}{n+1} = -\frac{n}{n+1}$$

$$C_n = 2(n+1) \left( H_n - \frac{n}{n+1} \right)$$

$$= 2(n+1)H_n - \frac{2n(n+1)}{n+1}$$

CF

$$C_n = 2(n+1)H_n - 2n, \quad C_0 = 0$$

$$C_0 = 0, \quad C_1 = 1, \quad C_2 = 2 \cdot 3 \cdot \frac{3}{2} - 4 = 5 \text{ etc.}$$

# GEOMETRIC SUM REVISITED

$$① \quad S_n = \sum_{k=0}^n ax^k$$

$$x^{k+1} = \overset{k}{\circlearrowleft} x \cdot \overset{k}{\circlearrowleft} x$$

x is a constant with respect to k

② Observe:

$$\sum_{k=0}^n ax^{k+1} = x \sum_{k=0}^n ax^k$$

## PERTURBATION TECHNIQUE

EVALUATE

$$S_n + ax^{n+1} = ax^0 + \sum_{k=0}^n ax^{k+1}$$

$S_{n+1}$

$$\stackrel{②}{=} a + x \sum_{k=0}^n ax^k$$

$$\stackrel{①}{=} a + xS_n$$

Solve for  $S_n$ :

$$S_n - xS_n = a - ax^{n+1}$$

$$S_n = \frac{a(1-x^{n+1})}{1-x}$$

# ANOTHER EXAMPLE OF PERTURBATION TECHNIQUE

EVALUATE

$$S_n = \sum_{k=0}^n k 2^k$$

$$2^{k+1} = 2 \cdot 2^k$$

$$S_n + (n+1)2^{n+1}$$

$$= \sum_{k=0}^n (k+1)2^{k+1}$$

$S_{n+1}$

$$= \sum_{k=0}^n k \cdot 2^{k+1} + \sum_{k=0}^n 2^{k+1}$$

geom

$$= 2 \sum_{k=0}^n k \cdot 2^k + 2^{n+2} - 2$$

$$= 2S_n + 2^{n+2} - 2$$

EVALUATE  $S_n$

$$S_n + (n+1)2^{n+1} = 2S_n + 2^{n+2} - 2$$

$$S_n(1-2) = -(n+1)2^{n+1} + 2^{n+2} - 2 \quad | \cdot (-1)$$

$$S_n = 2^{n+1}(n+1-2) + 2$$

$$S_n = (n-1)2^{n+1} + 2 \quad \text{CLOSED FORMULA}$$

$$\sum_{k=0}^n k 2^k = (n-1)2^{n+1} + 2$$

PUT  $2 = x$  and REPEAT

$$\sum_{k=0}^n k x^k = \frac{x - (n+1)x^{n+1} + nx^{n+2}}{(1-x)^2}$$

Use  $\sum_{k=0}^n x^{k+1} = \frac{1-x^{n+2}}{(1-x)}$

HOME.